

# Solidity Optimizer

## Solidity Summit

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Which one is more gas efficient<sup>1</sup>? a() or b()?

```
contract C {
    uint[] arr = [...];

    function a() external returns (uint sum) {
        for(uint i = 0; i < arr.length; i++) {
            sum += arr[i];
        }
    }

    function b() external returns (uint sum) {
        uint[] memory arr_copy = arr;
        uint length = arr_copy.length;
        for(uint i = 0; i < length; i++) {
            sum += arr_copy[i];
        }
    }
}
```

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<sup>1</sup>Inspired by a post from Patrick Collins.

# How does the compiler work?

1. Solidity code is parsed into an Abstract Syntax Tree (AST).
2. Perform analysis on the AST.
3. Code generation:
  - 3.1 Legacy code generation: translate AST into EVM bytecode directly<sup>2</sup>.
    - ▶ Perform bytecode based optimizations.
  - 3.2 Intermediate representation: translate AST into Yul.  
(`--via-ir` or `viaIR: true`)
    - ▶ Perform Yul optimizations.
    - ▶ Translate Yul into EVM bytecode.
    - ▶ Perform bytecode based optimizations.

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<sup>2</sup>Yul and its optimizer is partially used.

## Bytecode based optimizer

- ▶ Usually works across basic blocks.
- ▶ Simple evaluation of expressions.
  - ▶  $\text{add}(A, B) \rightarrow A + B$
  - ▶ Find more on RuleList.
- ▶ Cannot perform complex optimizations.

## Where does the bytecode based optimizer fail?

- ▶ Rule<sup>3</sup>:  $\text{mul}(a, 2) \rightarrow \text{shl}(1, a)$ .

```
function f(uint256 a) public pure returns (uint256) {
    unchecked {
        return a * 2;
    }
}
```

- ▶ Basic block: JUMPDEST, PUSH 2, MUL.
- ▶ Because the value a is outside the basic block, this rule cannot be applied.
- ▶ Engineering decision: we want to keep the bytecode based optimizer as simple as possible.

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<sup>3</sup>Based on a question by Alexey.

# Yul Optimizer

- ▶ Can perform more complex optimizations across blocks.

```
function f(uint256 a) public pure returns (uint256) {
    unchecked {
        return a * 2;
    }
}
```

Approximate IR:

```
function f(a) -> r {
    r := mul(x, 2)
}
```

Optimized into:

```
function f(a) -> r {
    r := shl(1, x)
}
```

- ▶ The `--via-ir` codegen can optimize `mul(x, 2)` into `shl(x, 1)`.

## What can Yul do better?

- ▶ Remove redundant division by zero checks.
- ▶ More inlining.
- ▶ Better stack management.
- ▶ Packed structs are better optimized.
- ▶ Trivial things like "`x != 0` instead of `x > 0`".
- ▶ Small and independent steps that can be executed in a sequence.

## Loop invariant code motion

```
function a() external returns (uint sum) {
    for(uint i = 0; i < arr.length; i++) {
        sum += arr[i];
    }
}
```

Translated Yul excluding a certain optimization:

```
let sum := 0
for {let i := 0 } lt(i, sload(0)) {
    // overflow check for i
    if eq(i, not(0)) { panic() }
    i := add(i, 1)
}
{
    let arr_value_i := sload(add(HASH, i))
    // Overflow check for +=
    if gt(arr_value_i, not(sum)) { panic() }

    sum := add(sum, arr_value_i)
}
```



## Loop invariant code motion

The length (`sload(0)`) is an invariant in the loop!

Optimized code:

```
let sum := 0
let len := sload(0)
for {let i := 0 } lt(i, len) {
  // overflow check for i
  if eq(i, not(0)) { panic() }
  i := add(i, 1)
}
{
  let arr_value_i := sload(add(HASH, i))
  // Overflow check for +=
  if gt(arr_value_i, not(sum)) { panic() }

  sum := add(sum, arr_value_i)
}
```

## The other way to compute sum

Copying array into memory first. Then do all computations there.  
Has memory overhead.

```
function b() external returns (uint sum) {  
    uint[] memory arr_copy = arr;  
    uint length = arr_copy.length;  
    for(uint i = 0; i < length; i++) {  
        sum += arr_copy[i];  
    }  
}
```

In legacy codegen, this is cheaper. But viaIR makes the simple implementation cheaper!

## A note on benchmarking

- ▶ Benchmark optimizations.
- ▶ Function dispatch can affect gas.
- ▶ Look at diffs of assembly as well as IR.

## Dispatch affecting gas

```
contract C {
    function a() external {}
    function b() external {}
    function c() external {}
}

contract CByteCode {
    // pseudocode for function dispatch of C
    fallback() external {
        if (msg.sig == 0x0dbe671f)
            a();
        else if (msg.sig == 0x4df7e3d0)
            b();
        else if (msg.sig == 0xc3da42b8)
            c();
    }
    function a() internal {}
    function b() internal {}
    function c() internal {}
}
```

## Dispatch affecting gas

- ▶ The order of the function in the dispatch can introduce an overhead.
- ▶ For benchmarking gas, try to have contracts with only a single function.
- ▶ Certain frameworks can display the gas without the dispatch overhead.

## Diffs of assembly or IR

- ▶ The `--asm` option is more readable than the binary.
- ▶ The `--ir-optimized --optimize` is even more readable.
- ▶ [godbolt.org](https://godbolt.org) now has Solidity support!

## Future plans for the optimizer

- ▶ General idea: intuitive code is also optimal.
- ▶ Improving inlining heuristics.
- ▶ Complex optimizations that require symbolic reasoning:

```
function f() external returns (uint sum) {  
    for (uint i = 0; i < arr.length; i++) {  
        sum += arr[i];  
    }  
}
```

- ▶ Remove redundant reverts.
- ▶ Optimize memory usage.

## Safety of the optimizer

- ▶ General fuzzing and differential fuzzing by Bhargava.
- ▶ External fuzzing efforts from academia (Alex Groce, etc).
- ▶ Formally verify some steps in Z3.
- ▶ Symbolically check some optimization steps from dapptools.
- ▶ Hand written proofs, PR reviews, etc.

No known compiler bug on a deployed contract!



# Slides

<https://hrkrshnn.com/t/devconnect.pdf>